CSE11 Lecture 20 Fall 2013 Recursion

#### Recursion

- **recursion**: The definition of an operation in terms of itself.
  - Solving a problem using recursion depends on solving smaller or simpler occurrences of the same problem.
- **recursive programming**: Writing methods that call themselves to solve problems recursively.
  - An equally powerful substitute for *iteration* (loops)
  - Particularly well-suited to solving certain types of problems

#### Why learn recursion?

- "Cultural experience" think differently about problems
- Solves some problems more naturally than iteration
- Leads to elegant, simplistic, short code (when used well)
- Many programming languages ("functional" languages such as Scheme, ML, and Haskell) use recursion exclusively (no loops)

#### Simple Exercise

How many students total are directly to the left of you in your "row" of the classroom?

- You can only see the person right next to you
- But, You can ask that person a question and he/she can respond to you
- How can we solve this problem (recursively)?



#### **Recursive algorithm**

- Number of people to my left
  - If there is someone to my left, ask him/her how many people are to their left
    - When they respond with a value N, then I will answer N + 1.
- If there is nobody to my left, I will answer **0**. Call Stack You Person next to You

#### Recursion and cases

- Every recursive algorithm involves at least 2 cases:
  - **base case**: A simple occurrence that can be answered directly.
  - recursive case: A more complex occurrence of the problem that cannot be directly answered, but can instead be described in terms of smaller occurrences of the same problem.
  - Some recursive algorithms have more than one base or recursive case, but all have at least one of each
  - A crucial part of recursive programming is identifying these cases.
  - Can also create/define Recursive Structures.

#### **Complex objects**

 How might you design a class called NestedRects of graphical objects that look like this?



- Requirements for the constructor:
  - Like many graphical objects, takes 5 parameters:
    - x and y describing coordinates of upper left
    - width and height of outermost rectangle
    - canvas
  - Spacing between rectangles is 4 pixels

```
Constructor for NestedRects
public NestedRects (double x double y,
                    double width, double height,
                    DrawingCanvas canvas) {
  new FramedRect(x, y, width, height, canvas);
 while (width \geq 8 \&\& height \geq 8) {
     width = width - 8;
     height = height - 8;
     x = x + 4;
     y = y + 4;
     new FramedRect(x, y, width, height, canvas);
```

#### Making NestedRects Useful

- Say that we want NestedRects objects to behave much like other graphical objects?
- NestedRects class should define methods like – moveTo()
  - removeFromCanvas()

But our constructor just draws the object

Need a way to keep track of entire collection of nested rectangles

Could use arrays for an iterative solution, but lets pretend we don't know about arrays.

#### Challenges

- Need to keep track of the rectangles in the collection
- Instance variables for each of the rectangles won't work:

FramedRect rectangle1, rectangle2;

We don't know how many there will be until a user specifies parameters when constructing one

#### A Recursive Solution

- A recursive structure consists of
  - A base structure (the simplest form of the structure)
  - A way to describe complex structures in terms of simpler structures of the same kind
- Let's change the way we think about NestedRects
  - Rather than a series of FramedRects...
  - = outer FramedRect + a smaller NestedRects inside



#### NestedRects: a recursive def'n

public class NestedRects {

private FramedRect outerRect; // outermost rectangle
private NestedRects rest; // inner nested rects

```
// Move nested rects to (x, y)
public void moveTo(double x, double y) {
    outerRect.moveTo(x, y);
    if (rest != null) {
        rest.moveTo(x+4, y+4);
     }
}
```

```
// Remove the nested rects from the canvas
public void removeFromCanvas() {
    outerRect.removeFromCanvas();
    if (rest != null) {
        rest.removeFromCanvas();
    }
```

Tracing the execution of new NestedRects( 50, 50, 19, 21, canvas);







#### A Better Recursive Solution?

- moveTo and removeFromCanvas require checking whether rest is null
- Missing check will cause program to crash

• Can we write NestedRects to avoid the check for null?

#### Two Kinds of NestedRects

- "Normal" recursive case
  - outerRect
  - rest
- A special "simplest" NestedRects: empty!

Define a new class, BaseRects, representing an empty collection of FramedRects

### A Simple Base Class (an Empty Nested Rect)

public class BaseRects extends NestedRects2 {
 // Constructor has nothing to initialize
 public BaseRects() { }

// Move nested rectangles to (x, y)
public void moveTo(double x, double y) { }

// Remove nested rectangles from canvas
public void removeFromCanvas() { }

### A Base Class (an Empty Nested Rect)

public class BaseRects extends NestedRects2 {
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 public BaseRects() { }

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// Remove nested rectangles from canvas
public void removeFromCanvas() { }

#### **Revised Recursive Class**

public class NestedRects2 { private FramedRect outerRect; private NestedRects2 rest; // inner nested rects

// outermost rectangle

```
public NestedRects2(double x, double y,
                   double width, double height,
                   DrawingCanvas canvas) {
     outerRect = new FramedRect(x, y, width, height, canvas);
     if (width \geq 8 \&\& height \geq 8) {
             rest = new NestedRects(x+4, y+4, width-8,
                                      height-8, canvas);
    } else { // construct a base object
             rest = new BaseRects();
```

// Move nested rects to (x, y)
public void moveTo(double x, double y) {
 outerRect.moveTo(x, y);
 rest.moveTo(x+4, y+4)
}

// Remove the nested rects from the canvas
public void removeFromCanvas() {
 outerRect.removeFromCanvas();
 rest.removeFromCanvas();

## Evaluating new NestedRects2(54, 54, 11, 13, canvas)





 Since objects of type BaseRects and NestedRects2 know how to "moveTo" and "removeFromCanvas" ...

#### Checks for null are eliminated

'0' or Empty as the "base" case is often a good starting place for recursion

#### Designing recursive structures

Recursive structures built by defining classes for base and recursive cases

- Both implement same interface
- Base class
  - No instance variable has same type as interface or class
  - Generally easy to write
- Recursive class
  - At least one instance variable has same type as interface of class
  - Care needed to be sure methods terminate

# Recursive Methods (or Algorithms)

- Can write recursive methods that are not part of recursive structures
- SolveMe (N) --> X + SolveMe(N-1)
- A very common use of recursion are socalled "divide and conquer" algorithms
  - Solve two problems, each of 1/2 the size of the original, then assemble the full answer from both parts
  - Sorting in Searching (Chapter 20)

#### Base case replaces Base class

**Recursive methods** 

- Must include at least one base case
- Typically contain a conditional statement
  - At least one case is a recursive invocation
  - At least one case is a base case -- i.e., no recursive invocation of the method
- Without a BASE case you will recurse infinitely! (That's bad)

#### An example: Exponentiation

- Inspiration: Fast algorithms for exponentiation important to RSA algorithm for public key cryptography – calculate: B<sup>k</sup>
- A simple (not fast!) recursive method:

// returns base raised to exponent as long as exponent >=0
public double simplePower(double base, int exponent) {

```
if (exponent == 0) {
```

return 1;

```
} else {
```

return base \* simplePower(base, exponent-1);

#### An example: Exponentiation

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#### Rules for writing recursive methods

- Write the base case
  - No recursive call
- Write the recursive case
  - All recursive calls should go to simpler cases
  - Simpler cases must eventually reach base case

#### Applying rules to simplePower

- Base case: exponent == 0
  - Returns 1
  - Correct answer for raising base to the 0th power
  - No recursive invocation
- Recursive case: uses else clause
  - Recursive call involves smaller value for exponent
  - Recursive calls eventually reach base case of 0: exponent greater than 0 to start and always goes down by 1

#### Rules of Exponents

- Simple algorithm took advantage of these rules:
  - $-base^{0} = 1$
  - $-base^{exp+1} = base * base^{exp}$
- New algorithm will make use of this rule:
   Base<sup>m\*n</sup> = (base<sup>m</sup>)<sup>n</sup>

Let m = 2,  $n = \exp/2$ base<sup>exp</sup>= (base<sup>2</sup>)<sup>exp/2</sup>

#### Faster exponentiation

public double fastPower(double base, int exponent) {
 if (exponent == 0) {
 return 1;
 } else if ( exponent%2 == 1 ) { // odd exponent
 return base \* fastPower(base, exponent-1);
 } else {
 return fastPower(base\* base, exponent/2);
 }

base<sup>exp</sup>= (base<sup>2</sup>)<sup>exp/2</sup>

#### Tracing fastPower

fastPower(3, 16)

- = fastPower(9, 8)
- = fastPower(81, 4)
- = fastPower(6561, 2)
- = fastPower(43046721, 1)
- = 43046721 \* fastPower(43046721, 0)
- = 43046721 \* 1
- = 43046721

Only 5 multiplications!

Division by 2 is fast and easy for computers

#### Towers of Hanoi



Move all disks from left to right peg Move one at a time Can only put smaller disks on empty or larger disks

#### Base Case



Move all disks from left to right peg Move one at a time Can only put smaller disks on empty or larger disks

#### 1<sup>st</sup> Recursive Move



Move All Disks from Left Tower to Right Tower =

1

2

- Move 2,3,4 to Middle (Recursive) + Move 1 to Right (base)
- Then move Disks 2,3,4 (Recursive) to the Right Tower

## Suppose we've accomplished Step 1



After 1 has completed, then do 2 . Howework explores this. Note both of these are recursive.



To Move 2,3,4 to the Right Tower

2

= Move 3,4 to the left tower + move 2 the Right Tower



Moving 3,4 from middle tower to left tower



Can Now move 2 to the Right Tower



Move 3,4 to the Right Tower

- = Move 4 to middle + move 3 to the Right
  - Then move 4 to the right



#### Finishing



Move 4 to the Right (No recursion needed)



Move all disks from left to right peg Move one at a time Can only put smaller disks on empty or larger disks